

Things to know by May 2

Warning. This is not guaranteed to be everything you need to know.

The final exam is 9:00 AM on Wednesday, May 2, in Sitterson 011.

Know everything you were supposed to know for the two midterms and know something about the following terms:

base register	broadcast (message)	CSMA/CD
CSR	device driver	device registers
free token	Internet	limit register
local-area network	memory-mapped I/O	page fault
page table	physical memory	programmed I/O
switch table	token ring	virtual memory
virus		

The final is based on Chapter 1 through Subsection 8.1.4 (pp. 1-255), Chapter 10 up to Subsection 10.1.2.3 (pp. 312-322), and Sections 11.3 and 11.3 (pp. 382-388) of the textbook, the October 1987 *Scientific American* article “Networks for Advanced Computing” by Robert Kahn, the January 1987 *IEEE Network* article “The IBM Token-Ring Network: A Functional View” by Norman Strole, Section 5.5.3 (pp. 209-218) of *Computer Science: A Modern Introduction* by Goldschlager and Lister, and Chapter 2 (pp. 13-25) of *Internetworking with XINU* by Douglas Comer.

Understand what a device driver is, how it does its task, and how the operating system is interfaced to it.

Understand how the operating system manages page tables.

Understand how token rings and Ethernets allocate the communication channel.

In addition to those problems you looked at the first two midterms, you might consider exercises 1 and 11 of Chapter 10 of Bach’s book and exercises 2.9 and 2.10 (p. 25) of Comer’s.

The last five pages of this handout is a *copy* of *last* year’s final exam in Comp 190. Last year we also covered windowing systems. One of the *short definitions* deals with that topic.

Final exam–May 2, **1989**
Closed book section (40 points)

The exam is to be turned in at 5:00 PM.

Work the closed book section first and turn it in before you consult your books and notes to work on the open book section. For the closed book section, write your answers on the exam itself. For the open book section, write your answers on separate pieces of paper.

University regulations require that you sign the following pledge on the first page of your turned-in exam.

I have neither received nor given any unauthorized aid on this exam. _____

Problem 1. (24 points–3 points each)

Give short definitions (one or two phrases or sentences) of the following terms.

datagram

hashing

inode

ports (concerning sockets)

process table

switch table

Unix

window (as in the X window system)

Problem 2. (12 points–3 points each)

Give a brief description of what the following UNIX system calls do at the user level. *You don't need to describe the implementation but say something about what they return.*

`execve(path, argv, envp)`

`fork()`

`sendto(s, msg, msg_len, flags, to, to_len)`
ignore the msg_len, flags, and to_len parameters

`setuid(uid)`
be sure to mention when it fails

Problem 3. (4 points)

Name the three regions of a Unix process and describe *briefly* (five words or less for each) their contents.

Final exam—May 2, **1989**

Open book section (60 points)

The exam is to be turned in at 5:00 PM. The closed book section should be turned in before you open your books and notes to work the open book section. For the open book section, write your answers on separate pieces of paper.

Problem 1: (5 points)

Suppose a group called **stooges** has been created with users **moe**, **curley**, and **larry**.

Subproblem 1.a: (1 point)

Can **curley** create a file which is readable only by **stooges**? How?

Subproblem 1.b: (4 points)

Suppose **moe** wants to create a program **nyack** that allows either **curley** or **larry** to send signals to any process **moe** is running as if **moe** sent the signal himself. Can **moe** (assuming he has worked hard in Comp 190) write this program? Can he restrict **nyack** so that only **stooges** can run it?

Problem 2: (5 points)

A process checks for signals when it enters or leaves the sleep state (if it sleeps at an interruptible priority) and when it returns to user mode from the kernel after completion of a system call or after handling an interrupt. Why does the process not have to check for signals when entering the system for execution of a system call?

This problem is exercise 12 on page 241 of Bach's book.

Problem 3: (6 points)

A networking hacker, having just studied the Unix code for obtaining locked inodes (Figure 4.3, page 64 of Bach), proclaims: "It's just like CSMA/CD." Questioned further he mutters something about sleeping on a locked inode and then checking the lock when awakened.

What do think is on the mind of this displaced communications geek?

Problem 4: (5 points)

We are snooping on a process running on a machine where the page size is 1000 bytes. We notice the following virtual to read address mappings: virtual 115 to real 19115, virtual 1492 to real 47492, virtual 1989 to real 47989, virtual 2338 to real 38338, and virtual 3456 to real 17456. We then observe that accessing virtual address 4545 results in a page fault.

Draw a page table for the first five pages used by this process.

Problem 5: (10 points)

You have been requested to write a someone unusual (and perhaps silly) device driver that will be used to determine whether or not some machine is transmitting on the Ethernet. Your device driver should behave as follows:

- (1) No matter how many bytes you try to read from this device, it will never actually read more than one byte. Thus “`read(fd, buff, 1000)`” and “`read(fd, buff, 1)`” both only change the first byte of `buff`.
- (2) The read byte is ‘B’ if the Ethernet is busy and ‘F’ if it is not.
- (3) All writes to this device have no effect. They are all successful, but any written data is ignored. (Writes to this device are just like writes to `/dev/null`.)
- (4) Your machine has a DEQNA ethernet interface like the one described in class and in Comer’s book. The only facts you need to know about the DEQNA is that its control and status register is stored at bus location 0174456 and that bit 13 of the CSR is 1 if a carrier is present on the ethernet and bit 12 of the CSR is 1 if the fuse is ok.

Subproblem 5.a: (7 points)

Outline with precision and brevity, how you will modify the kernel to implement this device driver.

Subproblem 5.b: (3 points)

Suppose the device driver is to be accessed via the special file `/dev/etherstat`. Describe how you must modify the file system to make this device accessible to user programs.

Problem 6: (11 points)

In Unix it is possible to mount a file system read-only by specifying the “`ro`” option to the “`mount`” command. Suppose a read-only file system `/dev/cd0` that has been mounted on directory `/nowrite` and that the root directory of this mounted file system contains a subdirectory `mtpnt`. Thus `/nowrite/mtpnt` is now directory within the read-only file system.

Subproblem 6.a: (2 points)

Attempts by the superuser to change the access permissions of `/nowrite/mtpnt` fail. Why?

Subproblem 6.b: (6 points)

Now, suppose the superuser wishes to mount a new file system `/dev/dsk5` on the directory `/nowrite/mtpnt`. The manual page for the `mount` system call does not seem to prohibit the use of a read-only directory as a mount point. That is, mounting on top of directories of read-only file systems is not listed as an error. How can the operating system perform a mount on a read-only file system?

Subproblem 6.c: (3 points)

Assuming the mount succeeds, now what happens if the superuser tries to change the access permissions of `/nowrite/mtpnt`?

Problem 7: (18 points)

The following program has been successfully compiled and run 100 times on a Unix system. It wrote “AA” to standard output 47 times, it wrote “AB” once, it wrote “ABA” 39 times, it wrote “ABB” 12 times, and it wrote “BAA” once.

Subproblem 7.a: (3 points)

Draw a picture showing the pipes created by the processes.

Subproblem 7.b: (2 points)

Why is two the least number of characters that can be written to standard output?

Subproblem 7.c: (2 points)

Why is three the greatest number of characters that can be written?

Subproblem 7.d: (3 points)

Describe how the program can produce “AA” as its output.

Subproblem 7.e: (4 points)

Describe how the program can produce “AB” as its output.

Subproblem 7.f: (4 points)

Describe how the program can produce “BAA” as its output.

```
#include <signal.h>
/*  1 */  main(argc, argv)
/*  2 */      int argc;
/*  3 */      char *argv[];
/*  4 */  {
/*  5 */      int p[2], pid;
/*  6 */      char b = 'X';
/*  7 */      pipe(p);
/*  8 */      write(p[1], "A", 1);
/*  9 */      pid = fork();
/* 10 */      if (pid > 0) {
/* 11 */          close(p[1]);
/* 12 */          read(p[0], &b, 1);
/* 13 */          write(1, &b, 1);
/* 14 */          kill(pid, SIGBUS);
/* 15 */          read(p[0], &b, 1);
/* 16 */          write(1, &b, 1);
/* 17 */          wait(0); }
/* 18 */      else {
/* 19 */          write(p[1], "B", 1);
/* 20 */          read(p[0], &b, 1);
/* 21 */          write(1, &b, 1);
/* 22 */          exit(0); }}
```