

**Problem 1 (5 points):**

Suppose a 10Mbps point-to-point link is set up between Asheville and Knoxville. The length of the link is about 120 km and data travels across the link at about  $2 \times 10^8$  m/s. What is the delay  $\times$  bandwidth product for the link?

**One way solution**

$$\begin{aligned} \text{Delay} &= 120 \text{ km} / (2 \times 10^8 \text{ m/s}) \\ &= 12 \times 10^4 \text{ m} / (2 \times 10^8 \text{ m/s}) \\ &= 6 \times 10^{-4} \text{ s} \\ &= .0006 \text{ s} \end{aligned}$$

$$\begin{aligned} \text{Delay} \times \text{Bandwidth} &= .0006 \text{ s} * 10 \text{ Mbps} \\ &= 6 \times 10^{-4} \text{ s} * 10^7 \text{ b/s} \\ &= 9 \times 10^3 \text{ b} \\ &= 6000 \text{ bits} \end{aligned}$$

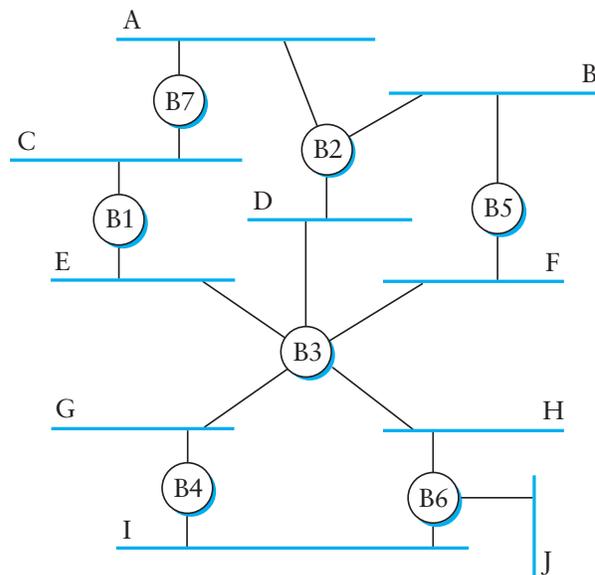
For round trip solution, double the delay and the delay  $\times$  bandwidth product

**Problem 2 (5 points)**

Given the extended LAN shown below, assume that bridge B1 and B5 suffer a near-simultaneous failure. Indicate which ports are **not** selected by the spanning tree algorithm after the recovery process and a new tree has been formed.

Clearly, the ports to B1 and B5 will not be used.

To avoid the cycle between the bridges and LANS, B3  $\rightarrow$  G  $\rightarrow$  B4  $\rightarrow$  I  $\rightarrow$  B6  $\rightarrow$  H  $\rightarrow$  B3, one of these ports must be disabled. The spanning tree algorithm will disable the port from B6 to I.



**Problem 3 (5 points)**

Suppose a computer (in this case it's `bulldog.unca.edu`) has the following routing table

Network and mask length	Interface or next hop
152.18.16/20	Interface tu3
152.18.32/20	Interface tu2
152.18.64/24	Interface tu0
152.18.224.16/28	Gateway 152.18.64.33
Default	Gateway 152.18.64.254

The IP numbers that this computer uses on its three interfaces are 152.18.16.9, 152.18.32.9, and 152.18.64.9. These IP numbers will not be used in the problem.

The leftmost column of the following table, gives an IP address. In the right column, write the routing for these addresses. Two examples are already given in the table.

Destination IP address	Next routing
152.18.16.33	Interface tu3
18.181.0.33	Gateway 152.18.64.254
152.18.17.33	Interface tu3
152.18.64.33	Interface tu0
152.18.65.33	Gateway 152.18.64.254
152.18.224.11	Gateway 152.18.64.254
152.18.224.21	Gateway 152.18.64.33

**Problem 4 (5 points)**

Suppose TCP's adaptive retransmission mechanism is running with  $\delta = 1/8$ ,  $\mu = 1$ , and  $\phi = 4$ . Assume that at the present time, *EstimatedRTT* is 800 and the *Deviation* is 400. The next two packets arrive with a measured RTT of 2000. What is the *TimeOut* value after the second packet arrives?

$$\text{Difference} = \text{SampleRTT} - \text{EstimatedRTT}$$

$$\text{EstimatedRTT} = \text{EstimatedRTT} + (\delta \times \text{Difference})$$

$$\text{Deviation} = \text{Deviation} + \delta \times (|\text{Difference}| - \text{Deviation})$$

$$\text{TimeOut} = \mu \times \text{EstimatedRTT} + \phi \times \text{Deviation}$$

$$\text{Difference} = 2000 - 800 = 1200$$

$$\text{EstimatedRTT} = 800 + (.125 \times 1200) = 950$$

$$\text{Deviation} = 400 + .125 \times (|1200| - 400) = 500$$

$$\text{Difference} = 2000 - 950 = 1050$$

$$\text{EstimatedRTT} = 950 + (.125 \times 1050) = 1081.25$$

$$\text{Deviation} = 500 + .125 \times (|1050| - 500) = 568.75$$

$$\text{TimeOut} = 1 \times 1081.25 + 4 \times 568.75 = 3356.25$$